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Contributions to the Construction of Extensible Semantic Editors



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Doctoral Dissertation, 2012

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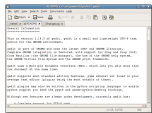


Orientation

Today's programming editors



Text editor



gedit, Notepad, ..

Semantic editor



Eclipse, NetBeans, ...

Background:

- Programming: From text to semantic editor
- Not all languages have semantic editors

Problems:

- Construction is time-consuming and complex
- Maintenance may be difficult (extensions)

Challenge:

How can we make it easier to construct
and maintain semantic editors?

Error feedback
Code browsing
Refactoring
Name completion

...



Approach

Generate services from specification

Specification: RAGs

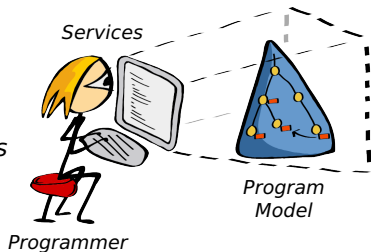
- Formalism, specify semantics
- Declarative, easily modularized
- JastAddJ, JModelica, JastAdd

RAGs – Reference Attribute Grammars

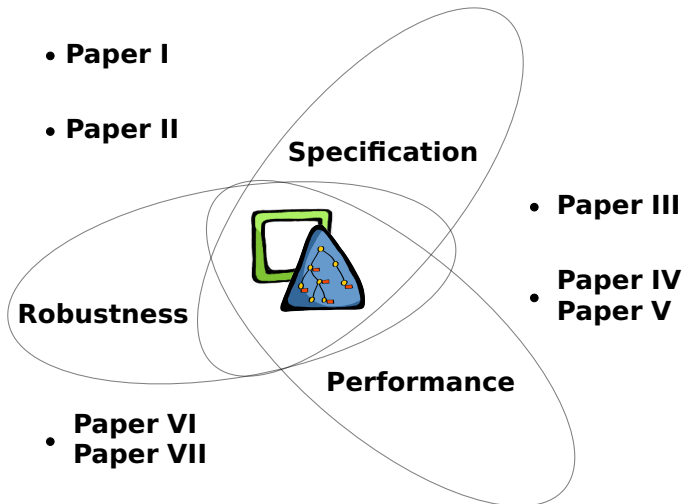
- *Grammar* – defines program model, abstract syntax tree (AST)
- *Attributes* – computed properties of AST nodes (types, scopes, ..)
- *References*: Attribute values

This dissertation:

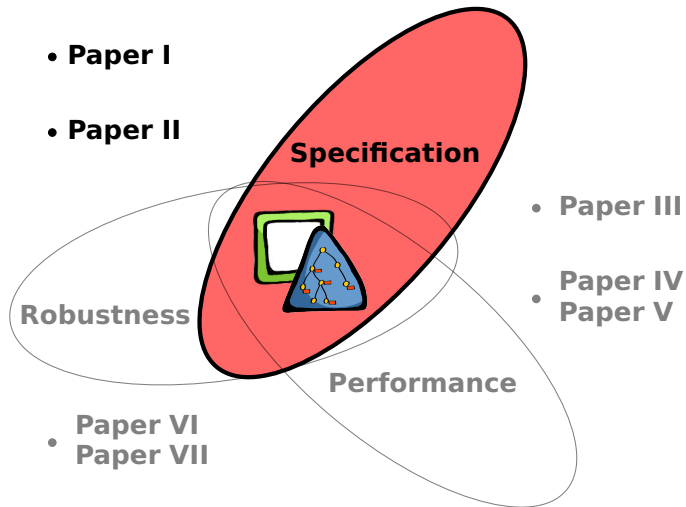
Extend compiler with editor module

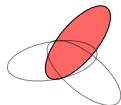


Contributions



Contributions





Specification

The JastAdd Editor Framework

6/20

Problem:

How to add an editor to an existing compiler?

Approach:

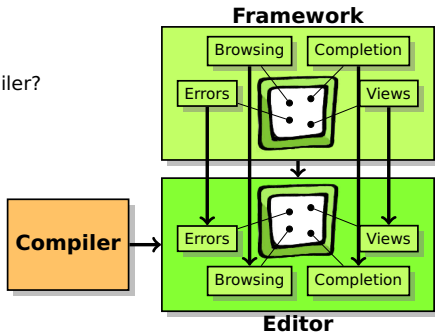
- Editor Framework
- JastAdd: semantics
- Eclipse: graphical components
- Predefined generic services
- RAG-based compiler extension

Results:

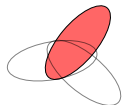
- Two demonstrators (size in LOC):
 - JastAdd: compiler 29,200, editor 4,300 (1,100)
 - PicoJava: compiler 210, editor 600 (420)

Conclusions

- Modularly defined editor services
- Compiler reuse (name analysis, type analysis, ...)



[Paper I]



Specification

Example Service:
Dead assignments

Problem:

How to specify flow analysis services?



```
a = 0;
b = 0;
a = b + 5;
a = a + b;
return a;
```

Here "a = 0" is a dead assignment because the value is not used and it could be removed.

Liveness (textbook)

Let n be a node and $\text{succ}[n]$ the set of successors for the node n :

$$\text{in}[n] = \text{use}[n] \cup (\text{out}[n] \setminus \text{def}[n])$$

$$\text{out}[n] = \bigcup_{s \in \text{succ}[n]} \text{in}[s]$$

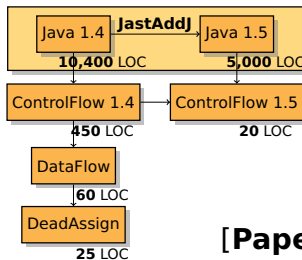
RAGs

```
syn Set CFGNode.in() circular [empty()] =
  use().union(out()).compl(def());
```

```
coll Set CFGNode.out() circular [empty()] with add;
Stmt contributes in() to CFGNode.out() for each pred();
Expr contributes in() to CFGNode.out() for each pred();
```

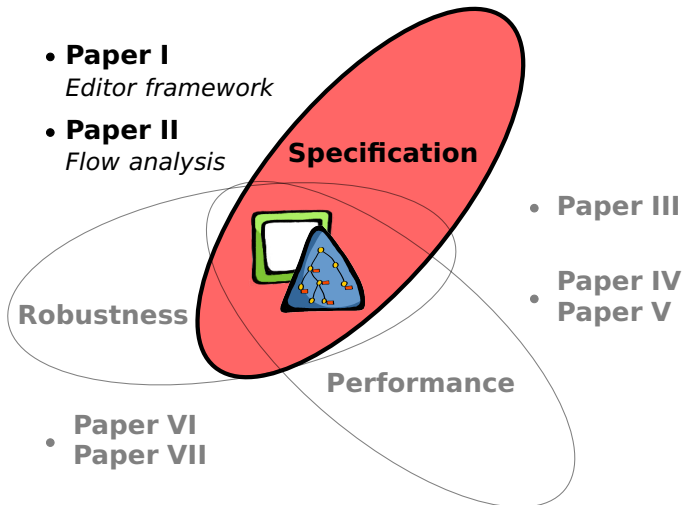
Conclusions:

- Textbook-like definitions
- Flow analysis added modularly with few LOC.
- Precision/performance on par with Soot.

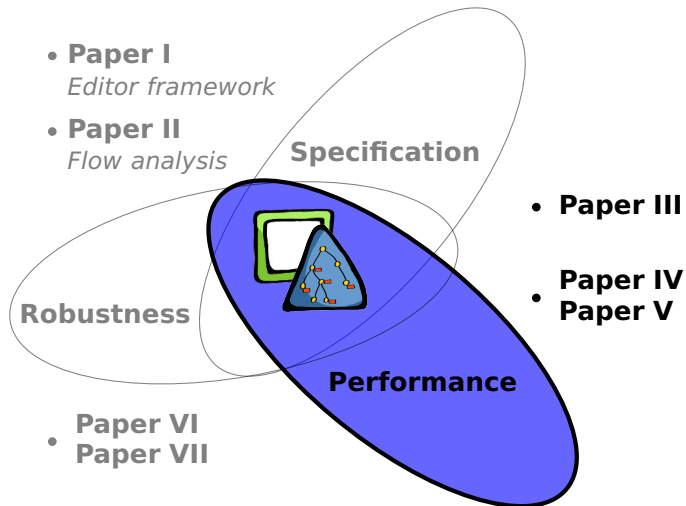


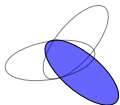
[Paper II]

Contributions



Contributions





Performance

Faster evaluation from scratch

Background: At attribute evaluation

- attribute dependencies \rightarrow call graph
- *no caching* - multiple evaluations \rightarrow *very slow*
- *full caching* - at most one evaluation \rightarrow *faster*

Problem: Memory/performance costs

New idea: Selective caching

- based on profiling
- skip caching of some attributes

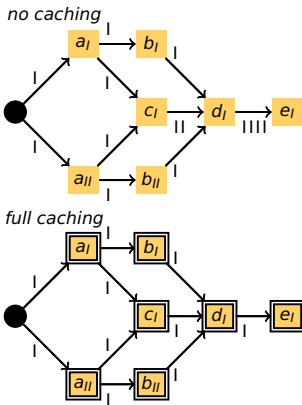
Results: 20% speedup and 38% memory reduction

- Compared to full caching
- JastAddJ Java compiler
- Java benchmarks

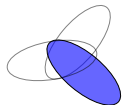
Reference Attributes

$a = b.c$ $b = d$ $e = ..$
 $c = d$ $d = e$

Attribute Instance Call Graphs



[Paper III]



Performance

Incremental evaluation

Problem:

How to efficiently update the program model after edits?

State of the art:

- Hand-coded solutions, complex, error-prone

Challenge:

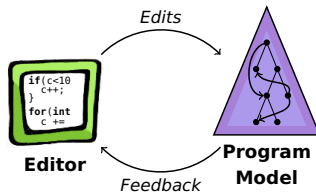
- Automatically update model, RAGs

Earlier work:

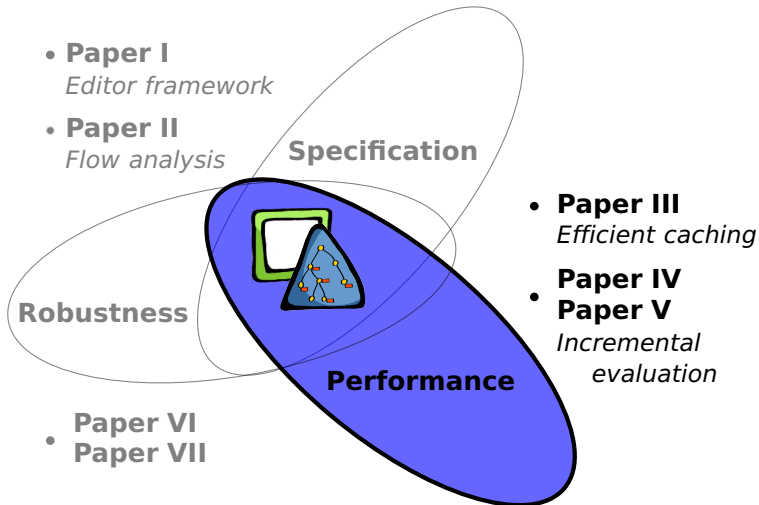
- Optimal automatic updates for AGs
- No handling of references

New results:

- Dynamic algorithm for RAGs.
 - Build dynamic dependency graph during evaluation
 - Use graph to uncache affected attributes after edits

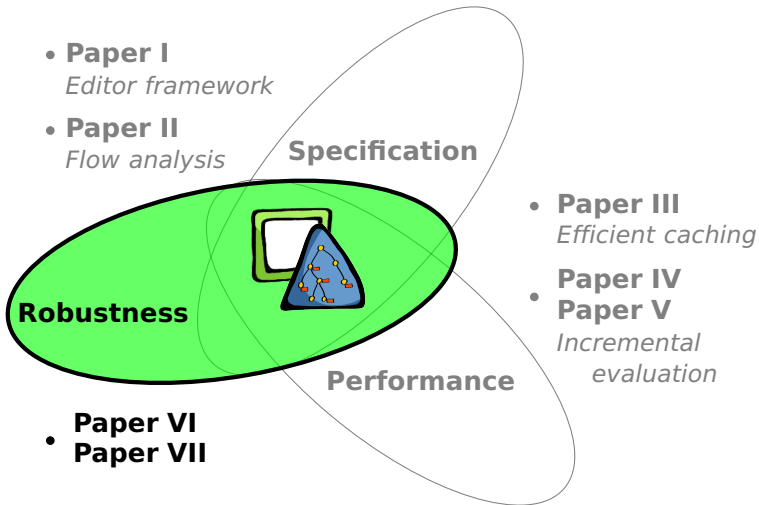


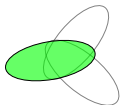
Contributions



Contributions

- **Paper I**
Editor framework
- **Paper II**
Flow analysis





Robustness

How to handle erroneous input?

Problem:

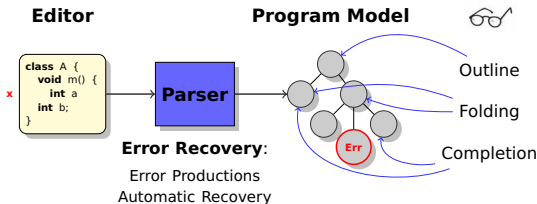
Scope errors cause recovery to fail

Idea:

- Use layout for recovery
- Aid existing recovery with preprocessor

New Algorithm:

Bridge parsing



```

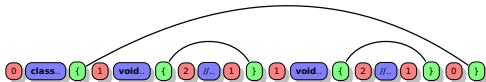
class C {
  void m() {
    // ..
  }
  void n() {
    // ..
  }
}
  
```

Island grammars:

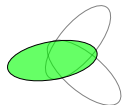
- Islands: Interesting ■
- Water: Uninteresting ■

Bridge Parsing

- Reefs: Patterns for recovery ■
- Bridges – Scopes

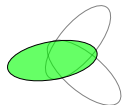


[Paper VI]



Robustness

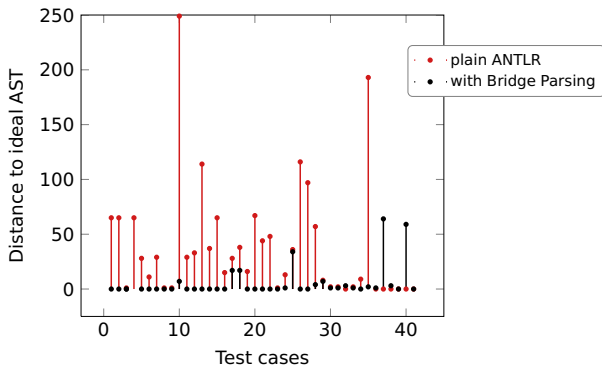
Bridge Parsing Algorithm



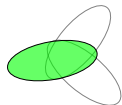
Robustness

Results from adding bridge parsing

Antlr – a well-known LL-based parser generator



[Paper VI]



Robustness

Bridge Parsing ideas in JSGLR

Collaboration: TU Delft

Problem:

Provide error recovery for Scannerless GLR (SGLR)

SGLR:

- *Generalized* LR: Arbitrary CFG
- *Scannerless* – include tokens in the grammar
- Language composition, e.g., Java-SQL and enum
- JSGLR – implementation of SGLR in Java

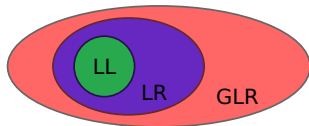
Method:

Recovery using island grammars,
layout and bridge parsing

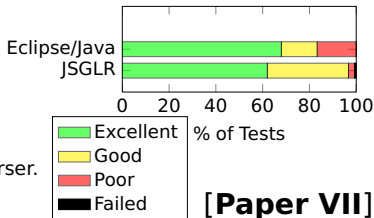
Results:

Recovery quality on par with the Eclipse Java parser.

Context Free Grammars



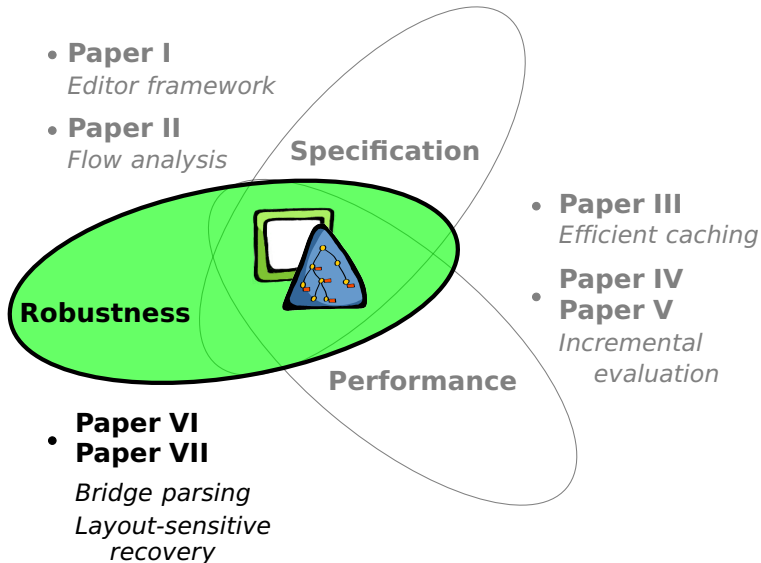
Recovery Quality



[Paper VII]

Bridge Parser Parts

Contributions



Contributions

