

**EDAA40**  
**Discrete Structures in Computer Science**

**9: A few words on proofs**

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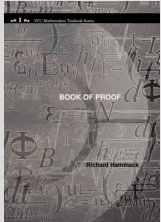
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This lecture is based on parts II and III of Richard Hammack's "Book of Proof".

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**definitions, theorems, proofs**

A **definition** is a statement that gives a precise meaning to a term or a symbol.

$$A \subseteq B \text{ iff } \forall x (x \in A \rightarrow x \in B)$$

$$n \in \mathbb{Z} \text{ is even iff } \exists k \in \mathbb{Z} (n = 2k)$$

$$n \in \mathbb{Z} \text{ is odd iff } \exists k \in \mathbb{Z} (n = 2k + 1)$$

A **theorem** is a statement that needs to be proven based on definitions (and axioms).

$$A \times (B \cap C) = A \times B \cap A \times C$$

$$\#(\mathbb{N}) < \#(2^{\mathbb{N}})$$

There are infinitely many prime numbers.

A **proof** is a chain of logical reasoning showing the truth of a theorem.

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## kinds of proofs

Proofs come in different flavors, which depend on the **form of the theorem**, and the chain of reasoning best suited to prove it.

Many theorems are conditional statements, i.e. they have the form "premise implies conclusion, or

$$P \Rightarrow C$$

$$\forall x \in \mathbb{Z} (x \text{ is odd} \rightarrow x^2 \text{ is odd})$$

$$\forall a, b, c \in \mathbb{Z} ((a|b \wedge b|c) \rightarrow a|c)$$

P	C	$P \Rightarrow C$
T	T	T
T	F	F
F	T	T
F	F	T

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## direct proof

**Theorem:** If P, then C.

**Proof:** Suppose P.

...

Therefore C.

**Theorem:**

$$x \text{ is odd} \rightarrow x^2 \text{ is odd}$$

**Proof:**

Suppose  $x$  is odd.

Therefore, there is an integer  $k$  such that  $x = 2k + 1$ .

$$\text{Thus } x^2 = (2k + 1)^2 = 4k^2 + 4k + 1 = 2(2k^2 + 2k) + 1.$$

Note that  $2k^2 + 2k$  is an integer.

Thus there is an integer  $n$  such that  $x^2 = 2n + 1$ .

Therefore  $x^2$  is odd.

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## direct proof with cases

Sometimes, the premise consists of several cases, and it becomes easier to study each case by itself.

$n$	$1 + (-1)^n(2n - 1)$
1	0
2	4
3	-4
4	8
5	-8
6	12

**Theorem:** If  $n \in \mathbb{N}$  then  $1 + (-1)^n(2n - 1)$  is a multiple of 4.

**Proof:** Suppose  $n \in \mathbb{N}$ . Then  $n$  is either even or odd.

**Case 1:** Suppose  $n$  is even. Then  $n = 2k$  for some  $k \in \mathbb{Z}$ .

$$\text{Thus } 1 + (-1)^{2k}(2(2k) - 1) = 1 + 1^k(4k - 1) = 4k.$$

That is a multiple of 4.

**Case 2:** Suppose  $n$  is odd. Then  $n = 2k + 1$  for some  $k \in \mathbb{Z}$ .

$$\text{Thus } 1 + (-1)^{2k+1}(2(2k+1) - 1) = 1 - (4k + 2 - 1) = -4k.$$

That is also a multiple of 4.

The result in both cases is a multiple of 4.

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## contrapositive proof

In some cases, it is easier to reason about a theorem in *contrapositive* form.

**Theorem:**  
If  $x^2 - 6x + 5$  is even, then  $x$  is odd.

**Proof:** Suppose  $x^2 - 6x + 5$  is even, i.e. there exists an integer  $a$  such that  $x^2 - 6x + 5 = 2a$ .

...

Thus there is an integer  $b$  such that  $x = 2b + 1$ .  
Therefore  $b$  is odd.

**direct proof:**

**Theorem:** If P, then C.

**Proof:** Suppose P.

...

Therefore C.

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## contrapositive proof

Contrapositive form:  $\neg C \Rightarrow \neg P$

P	C	$P \Rightarrow C$	$\neg C$	$\neg P$	$\neg C \Rightarrow \neg P$
T	T	T	F	F	T
T	F	F	T	F	F
F	T	T	F	T	T
F	F	T	T	T	T

**Theorem:** If  $x^2 - 6x + 5$  is even, then  $x$  is odd.

**Proof:**

Suppose  $x$  is even.

There is an integer  $a$  such that  $x = 2a$ .

$x^2 - 6x + 5 = 4a^2 - 12a + 4 + 1 = 2(2a^2 - 6a + 2) + 1$

So there is an integer  $b$  s.t.  $x^2 - 6x + 5 = 2b + 1$ .

Therefore  $x^2 - 6x + 5$  is not even.

**Theorem:** If P, then C.

**Proof:** Suppose not C.

...

Therefore not P.

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## proof by contradiction

Suppose we want to prove a proposition P, not necessarily in conditional form.

Proof by contradiction uses the fact that if we can show that not P results in a logical contradiction, e.g. it implies some conclusion C as well as its opposite, not C, then not P cannot be true, and so P must be true.

**Theorem:**  
If  $a, b \in \mathbb{Z}$  then  $a^2 - 4b \neq 2$ .

**Proof:** Suppose there are  $a, b \in \mathbb{Z}$  s.t.  $a^2 - 4b = 2$ .  
Since this implies  $a^2 = 4b + 2 = 2(2b + 1)$ ,  $a^2$  is even.  
Hence  $a$  is even, so  $a = 2c$  for some integer  $c$ .  
Thus  $4c^2 - 4b = 2$ , i.e.  $2c^2 - 2b = 1$ .  
Therefore  $2(c^2 - b) = 1$  with  $c^2 - b \in \mathbb{Z}$ .  
So 1 is even.

P	C	$\neg P$	$C \wedge \neg C$	$\neg P \Rightarrow C \wedge \neg C$
T	T	F	F	T
T	F	F	F	T
F	T	T	F	F
F	F	T	F	F

**Theorem:** P.

**Proof:** Suppose not P.

...

Or any other  
false proposition!  
Therefore C and not C.

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